

Graphs, trees and binary trees

1. Graphs
2. Trees
3. Binary trees

1. Graphs

A directed graph or **digraph** G is usually presented as a pair (V, E) :
 V is the **vertex set** and E is the **edge set**.

The elements of V are called **vertices**.

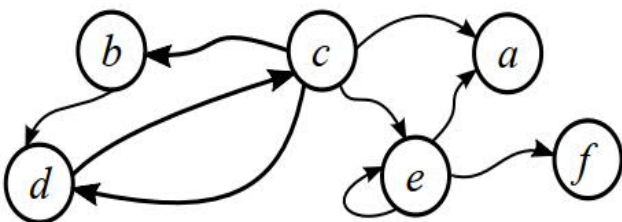
An element of E is called an **edge** and it is an **ordered pair** (a, b) of vertices.

For an ordered pair $(a, b) \neq (b, a)$.

Synonyms:

- vertex = node
- edge = arc = link

Example



$$V = \{a, b, c, d, e, f\}$$
$$E = \{(b, d), (d, c), (c, b), (c, d), (c, a), (c, e), (e, a), (e, e), (e, f)\}$$

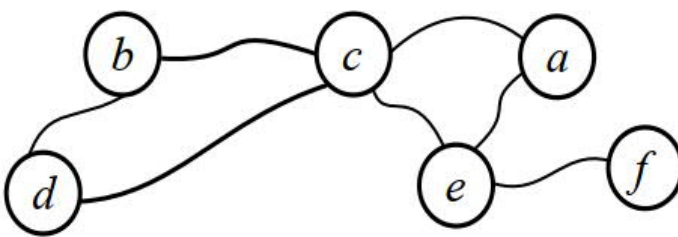
□

If the graph edges have no direction, they we have an **undirected graph**.

For an undirected graph $(a, b) = (b, a)$.

NOTE: multiple edges not allowed

Example



$$V = \{a, b, c, d, e, f\}$$

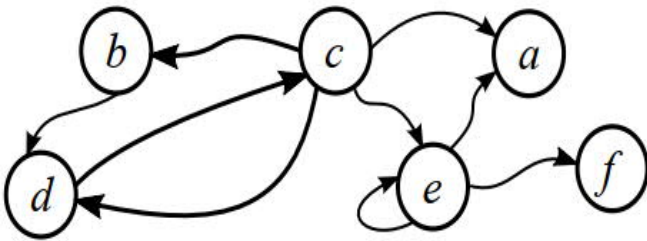
$$E = \{(b, d), (d, c), (c, b), (c, a), (c, e), (e, a), (e, f)\}$$

□

Terminology

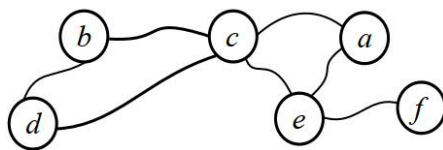
- If edge (a, b) exists in a digraph, then b is **adjacent** to a .
- If edge (a, b) exists in an undirected graph, then b is **adjacent** to a and a is **adjacent** to b .
- If edge (a, b) exists in a digraph, then a is **starting vertex** and b is the **final vertex**.
- In an undirected graph the **degree** $a \in V$ is the number of vertices that are adjacent to a .
- An **isolated** vertex in an undirected graph is one whose degree is 0.
- In a digraph, the **out-degree** of $a \in V$ is the number of edges leaving from a and the **in-degree** of $a \in V$ is the number of edges entering a .

Example



vertex x	vertices adjacent to x	vertices to which x is adjacent	in-degree	out-degree
a	—	c, e	2	0
b	d	c	1	1
c	b, d, e, a	d	1	4
d	c	b, c	2	1
e	e, a, f	c, e	2	3
f	—	e	1	0
□				

Example



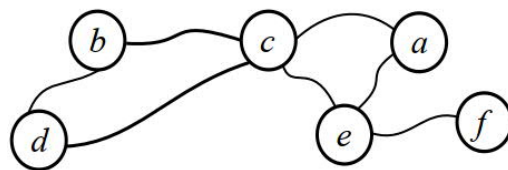
vertex x	a	b	c	d	e	f
vertices adjacent to x	c, e	d, c	a, e, b, d	b, c	c, a, f	e
degree	2	2	4	2	3	1

The graph has no isolated vertices. \square

More terminology

- A **path** of length k from vertex a_0 to vertex a_k is an ordered sequence of vertices $\langle a_0, a_1, a_2, \dots, a_k \rangle$ such that each edge (a_i, a_{i+1}) , $i = 0, 1, \dots, k - 1$ exists in the graph.
- In a **simple path** no vertex is repeated.
- If there is a path from a to b , then b is **reachable** from a .
- An undirected graph is **connected** if every vertex is reachable from every other vertex.
- A digraph is **strongly connected** if every vertex is reachable from every other vertex.
- The path $\langle a_0, a_1, a_2, \dots, a_k \rangle$ is a **cycle** when $a_0 = a_k$.
- In a **simple cycle** $\langle a_0, a_1, a_2, \dots, a_k \rangle$ the only repeated vertex is a_0 .
- An **acyclic** graph has no cycles.

Example



(i) Paths from b to a

$\langle b, c, a \rangle$ length = 2

$\langle b, d, c, e, a \rangle$ length = 4

$\langle b, d, c, b, d, c, a \rangle$ length = 6

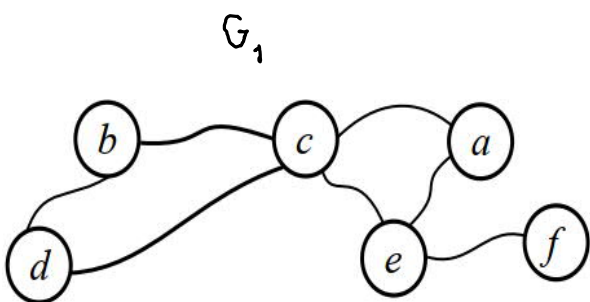
(ii) Simple paths from b to a
 $\langle b, c, a \rangle$ and $\langle b, d, c, e, a \rangle$

(iii) Every vertex is reachable from every other vertex. Hence the graph is connected.

(iv) Cycles in the graph
 $\langle b, d, c, b \rangle$ simple
 $\langle a, c, e, a \rangle$ simple
 $\langle a, c, b, d, c, e, a \rangle$ not simple

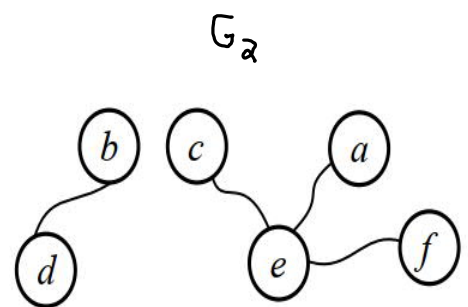
□

Example



G_1 is not acyclic.

□



G_2 is acyclic.

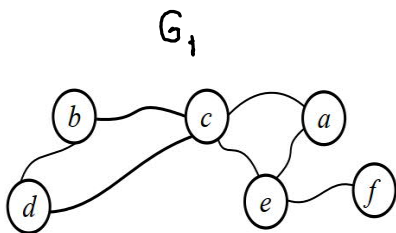
2. Trees

Undirected trees

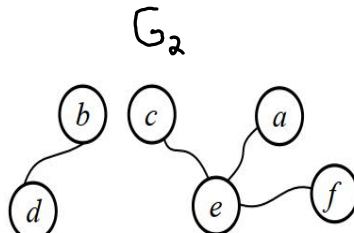
Each of the following defines an **undirected tree** $G = (V, E)$:

- G is connected and acyclic.
- G is connected and the number of edges is one less than the number of vertices.
- There is a unique simple path connecting every two vertices in G .
- G is acyclic, but adding any edge to E results in a graph with one cycle.

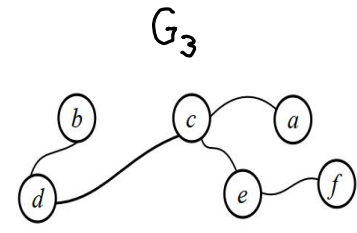
Example



not acyclic, so not a tree



no path from d to e ,
so not a tree



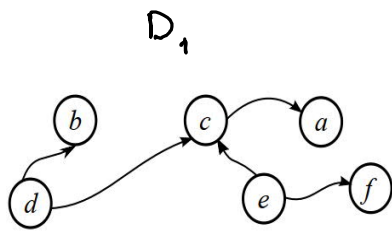
is a tree

□

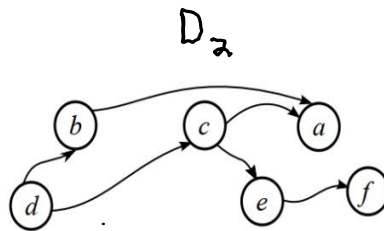
Rooted trees

A **rooted tree** is a digraph (V, E) where there is a unique simple path from one particular vertex, the **root** or r , to any other vertex, but there is no path from any vertex to r .

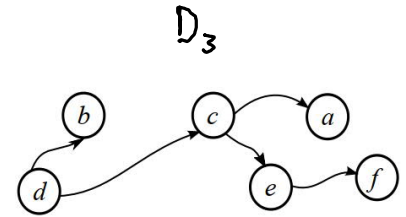
Example



no vertex has a path to all other vertices, so not a rooted tree



two paths from d to a , so not a rooted tree



is a rooted tree; root is d

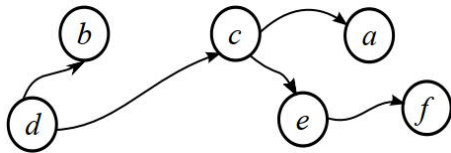


NOTE: usually omit arrows when drawing rooted tree and root is at the top

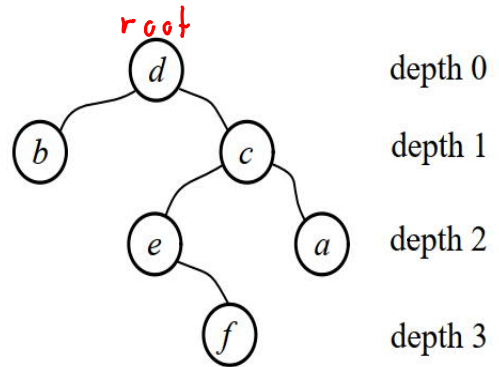
Terminology

- If a is a vertex on the unique path from r to b , then a is an **ancestor** of b .
- If a is an ancestor of b , then b is a **descendant** of a .
- The **subtree rooted at a** is a rooted tree having a as its root and includes all descendants of a .
- If (a, b) is an edge, then a is b 's **parent** and b is a ' **child**.
- Two vertices having the same parent are **siblings**.
- A vertex with no children is a **leaf** or an **external vertex**.
- A non-leaf vertex is an **internal vertex**.
- The **depth** of vertex a is the length of the simple path from r to a .
- The tree's **height** is largest depth of any vertex.

Example



OR



(i) path from d to f : $\langle d, c, e, f \rangle$

From this path we observe

- c is an ancestor of f
- f is a descendant of c

(ii) From edge (c, e)

c is e's parent and e is c's child

(iii) Vertices e and a are siblings.

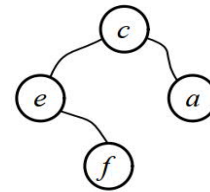
Vertices b and a are not siblings.

(iv) Vertices b, f and a are leaves.

Vertices d, e and c are internal vertices.

(v) The tree's height is 3.

(vi) The subtree rooted at c is



□

3. Binary trees

A **binary tree** is a rooted tree where every vertex has at most two children.

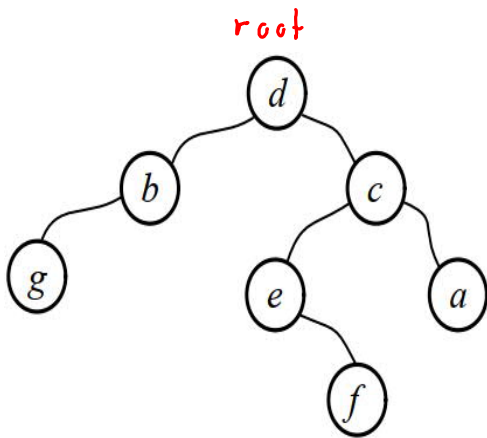
Terminology

- The **binary tree** with no vertices is the **empty tree** or the **null tree**.
- Each child of a vertex is either the **left child** or the **right child**.
- When a vertex has no left (right) child, then the left (right) child is **missing**.
- If b is the left (right) child of a , then the **left subtree** (**right subtree**) of a is the subtree rooted at b .

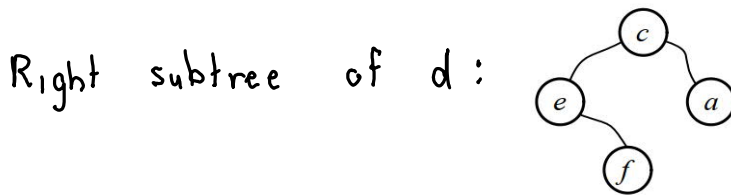
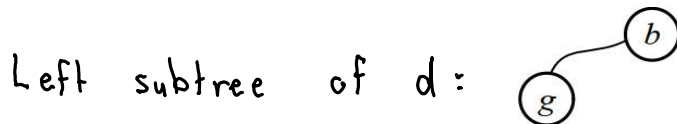
NOTE: both left and right subtrees are also binary trees

- In a **full** binary tree each vertex has either two children or no children.
- A **complete** binary tree is a full binary tree where all leaves are at the same depth.

Example

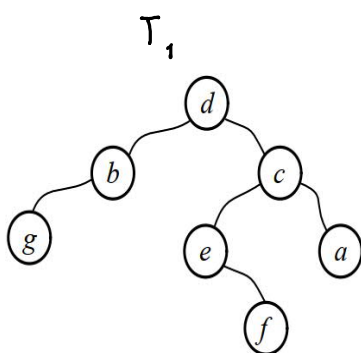


vertex	a	b	c	d	e	f	g
left child	-	g	e	b	-	-	-
right child	-	-	a	c	f	-	-
parent	c	d	d	-	c	e	b

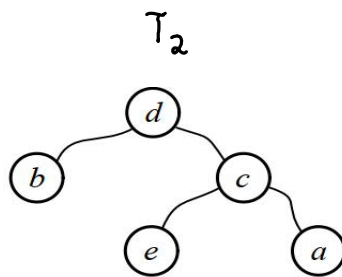


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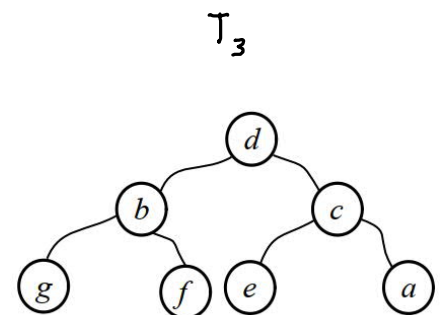
Example



b has only one child, so not a full binary tree



is a full binary tree, but not complete since leaves b and e are not at same depth



is a complete binary tree

□

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