

Mergesort: a divide and conquer sorting algorithm

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1. Introduction

Q: What does Mergesort do?

A: Mergesort will place elements in an array in order (smallest to largest or largest to smallest).

Mergesort properties:

- requires extra sorting space; to sort array $A[1..n]$ we need $Temp[1..n]$
- uses divide-and-conquer
- almost always presented as recursive
- uses **merging**

Q: What is merging?

A: Combining two sorted arrays into one sorted array.

Before merging: two sorted arrays $P[1..r]$ and $Q[1..s]$

$P[1]$	$P[2]$...	$P[r]$
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 $P[i] \leq P[i + 1]$ for all i

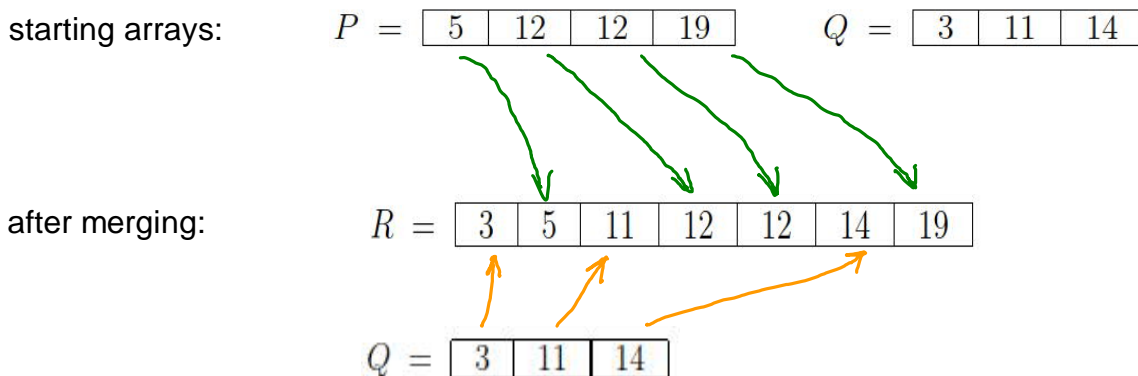
$Q[1]$	$Q[2]$...	$Q[s]$
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 $Q[i] \leq Q[i + 1]$ for all i

After merging: one sorted array that includes all elements from $P[1..r]$ and $Q[1..s]$

$$\boxed{R[1]} \quad \boxed{R[2]} \quad \dots \quad \boxed{R[r+s]} \quad R[i] \leq R[i+1] \text{ for all } i$$

Example



□

2. Merging

Merging sorted arrays $P[1..r]$ and $Q[1..s]$ into single sorted array $R[1..(r+s)]$ uses iterative strategy.

Iteration:

- add to end of R smallest element from either P or Q that has yet to be added, until either P or Q is empty

Final stage:

- add to end of R remaining elements of either P or Q , whichever is not empty

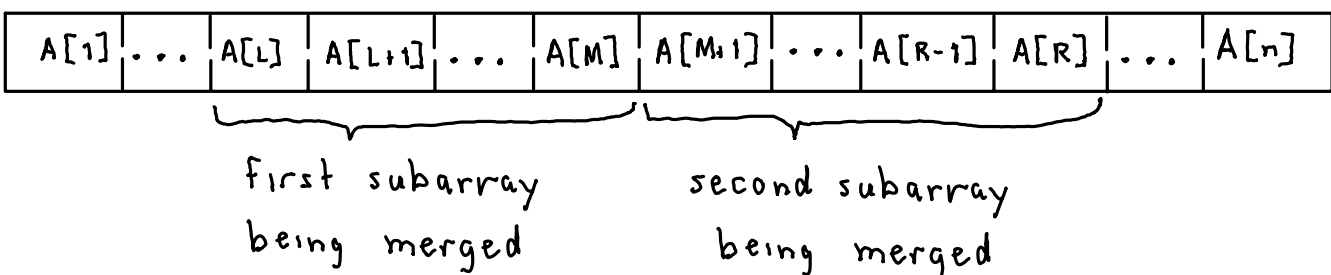
Pseudocode

```

1  MERGE(A, L, M, R)
2  input number array A, L is index of leftmost element to be
3  handled, R is index of rightmost element to be handled and M
4  is some index inbetween L and R,  $L \leq M \leq R$ 
5  /* On entering the procedure both subarray A[L..M] and subarray
6  A[(M+1)..L] are assumed to be sorted from smallest to largest.
7  On leaving the entire array A[L..R] is sorted from smallest to
8  largest. The temporary work array Temp[L..R] is used.*/
9  for i from L to R
10     Temp[i] = A[i]
11  end
12  /* Counter iL is the location in subarray Temp[L..M]. Counter iR
13  is the location in subarray Temp[(M+1)..R]. Counter iA is the
14  location in subarray A[L..R].*/
15  iL = L, iR = M + 1, iA = L
16  while iL ≤ M and iR ≤ R
17     if Temp[iL] ≤ Temp[iR] then
18         A[iA] = Temp[iL], iL = iL + 1
19     else
20         A[iA] = Temp[iR], iR = iR + 1
21     end
22     iA = iA + 1
23  end
24  /* We have either reached the end of subarray Temp[L..M] or the
25  end of subarray Temp[(M+1)..R]. Copy the remainder of the array
26  still containing elements to A. */
27  if iL > M then
28      ▷ copy Temp[iR..R] to A[iA..R]
29  else
30      ▷ copy Temp[iL..M] to A[iA..R]
31  end

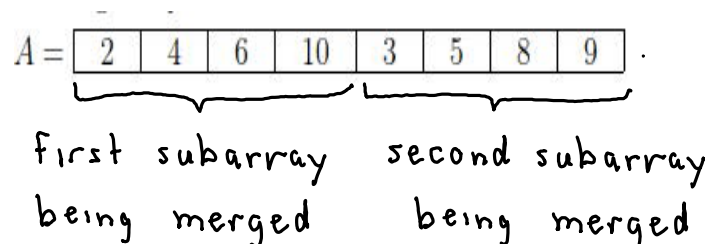
```

in MERGE



Example

starting array:



compute MERGE(A, 1, 4, 8)

after while-loop iteration iL iR iA arrays A and $Temp$ $\uparrow = iL$ $\uparrow = iR$ $\uparrow = iA$

0 1 5 1 $A =$

2	4	6	10	3	5	8	9
---	---	---	----	---	---	---	---

↑

$Temp =$

2	4	6	10	3	5	8	9
---	---	---	----	---	---	---	---

↑

↑

1 2 5 2

2	4	6	10	3	5	8	9
---	---	---	----	---	---	---	---

↑

2	4	6	10	3	5	8	9
---	---	---	----	---	---	---	---

↑

↑

2 2 6 3

2	3	6	10	3	5	8	9
---	---	---	----	---	---	---	---

↑

2	4	6	10	3	5	8	9
---	---	---	----	---	---	---	---

↑

↑

3 3 6 4

2	3	4	10	3	5	8	9
---	---	---	----	---	---	---	---

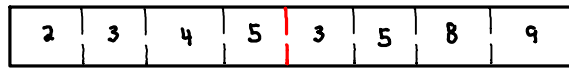
↑

2	4	6	10	3	5	8	9
---	---	---	----	---	---	---	---

↑

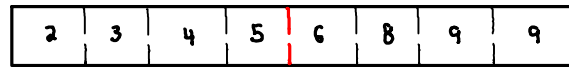
↑

4 3 7 5



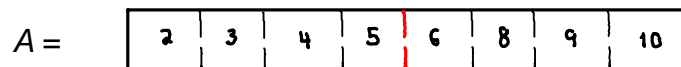
⋮

7 4 9 8



end of while-loop

After handling remainder of subarray $A[1..4]$ in line 30.



Comments on MERGE:

- to handle 'remainder' of one nonempty subarray **for**-loop is needed at line 28 and line 30
- 'remainder' of one nonempty subarray handled via test in while-statement (line 16) and lines 27-30; other approaches are possible
- merging can be done without work array Temp, but it is complicated
- similar approach can be used for merging arrays that have been rearranged according to other criteria

3. Mergesort computation

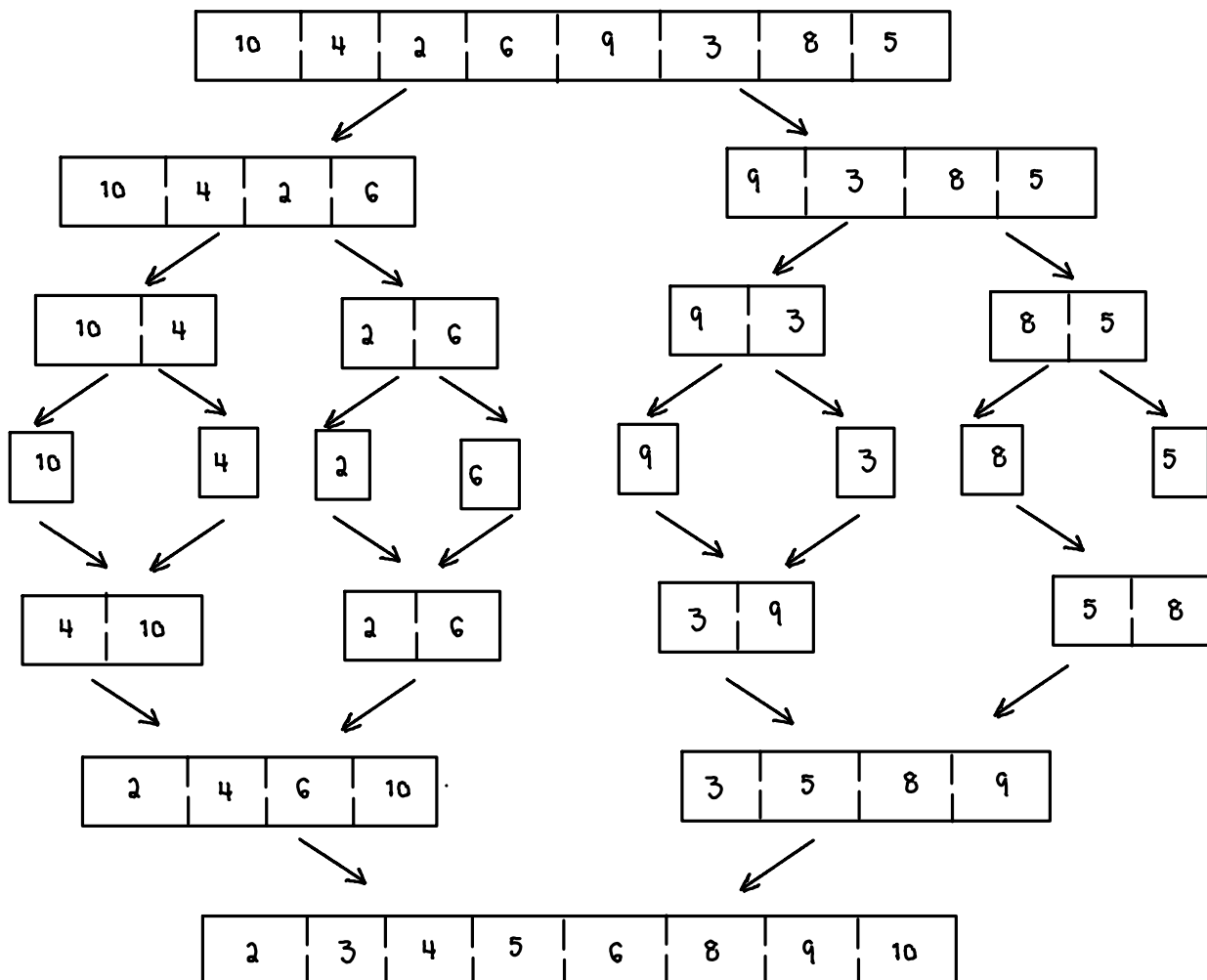
How is divide and conquer used in merge sort?

Divide: split starting array $A[1..n]$ in two equally sized halves: $A[1..M]$ and $A[(M+1)..n]$

Conquer: sort (recursively) both halves $A[1..M]$ and $A[(M + 1)..n]$

Combine: merge sorted halves $A[1..M]$ and $A[(M + 1)..n]$

Example



Recursion case: array to be sorted has at least 2 elements

Base cases: array to be sorted 1 or less elements

Mergesort pseudocode

```
1  MERGESORT(A, L, R)
2  input number array A, L is index of leftmost element to be
3  handled, R is index of rightmost element to be handled
4  /* We sort the subarray A[L..R] from smallest to largest using the
5  merge sort algorithm. */
6  if L < R then
7      M =  $\lfloor (L+R)/2 \rfloor$ 
8      MERGESORT(A, L, M)
9      MERGESORT(A, M+1, R)
10     MERGE(A, L, M, R)
11 end
```

Example

starting array: $A =$

10	4	2	6	9	3	8	5
----	---	---	---	---	---	---	---

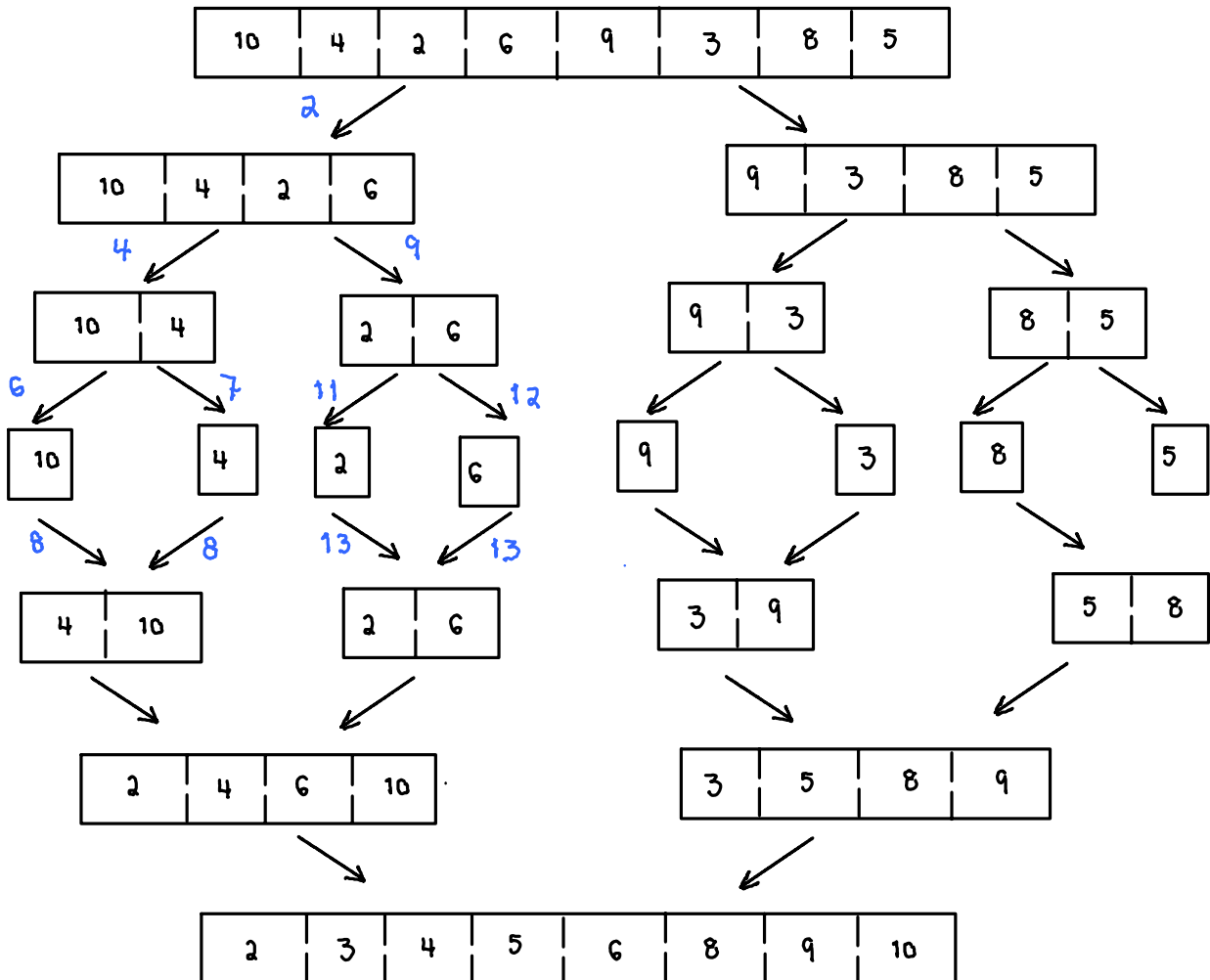
compute MERGESORT(*A*, 1, 8)

step	recursion level	code line(s)	computation	array A	↑ = L	↑ = R	↑ = M
1	1	7	$L=1, R=8, M=4$	10 4 2 6 9 3 8 5	↑		↑
2	1	8	MERGESORT(A, 1, 4)				
3	2	7	$L=1, R=4, M=2$	10 4 2 6 9 3 8 5	↑	↑	↑
4	2	8	MERGESORT(A, 1, 2)				
5	3	7	$L=1, R=2, M=1$	10 4 2 6 9 3 8 5	↑	↑	↑
6	3	8	MERGESORT(A, 1, 1)				
7	3	9	MERGESORT(A, 2, 2)				
8	3	10	MERGE(A, 1, 1, 2)	4 10 2 6 9 3 8 5			
9	2	9	MERGESORT(A, 3, 4)				
10	3	7	$L=3, R=4, M=3$	10 4 2 6 9 3 8 5	↑	↑	↑
11	3	8	MERGESORT(A, 3, 3)				
12	3	9	MERGESORT(A, 4, 4)				
13	3	10	MERGE(A, 3, 3, 4)	4 10 2 6 9 3 8 5			

□

Example (contd)

Order used in MERGESORT. *step value*



Comments on MERGESORT:

- amount of pseudocode deceptive since MERGE does all the work
- efficiency is not sensitive to starting order of array (unlike quicksort)
- can be implemented without recursion

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tekijä: Frank Cameron

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