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NOTE				
• individual node can be accessed by its key				
In C++ for general tree:				
1 struct Node		_		
2 { 3 // key = a unique identifier for the n	ode			
4 int key; 5 // name = name we give to the node				
<pre>6 std::string name; 7 // children = pointers to the children</pre>	of the node			
<pre>8 std::vector<node*> children; 9 // parent = pointer to the parent of t</node*></pre>	he node			
<pre>10 Node* parent; 11 }</pre>				
In C++ for binary tree:				
1 struct Node 2 {				
<pre>3 // key = a unique identifier for the r 4 int key;</pre>	iode —			
5 // name = name we give to the node 6 std::string name;	-			
7 // leftChild = pointer to the left chi 8 Node* leftChild; 9 // rightChild = pointer to the left chi				
10 Node* rightChild; 11 // parent = pointer to the Pert of the Per				
12 Node* parent; ————————————————————————————————————				

	ΓES																
- If r	ode d	oes no	ot exi	st, th	en po	inter	value	is n	ıllp	tr.							
- Inc	dividua	l node	is a	ccess	sed by	y its :	key.										
	node er conta			-								s prob	ably	ineffi	cient.	Some	Э
- All	tree n	odes	shoul	d be	store	d in a	a con	tainer	. For	exan	nple:						
				sto	l::uno	rdere	d_map	<int,< td=""><td>Node></td><td>allNo</td><td>des;</td><td></td><td></td><td></td><td></td><td></td><td></td></int,<>	Node>	allNo	des;						
3. Tr	avers	als															
Q: WI	nat is a	trave	rsal?)													
A: Wh	nen all	of the	elen	nents	of a	data	struc	ture a	re vis	ited.							
Single	dimer	nsiona	l arra	y tra	versa	ıl:											
	4	િ		<u>ځ</u>		1	B	~ E		3	9		7				
1																	
	consid	ler thr	ee tr	ae tra	avers	als: r	reord	er tra	versa	l nos	torde	r trav	ersal	and i	norde	r trav	ers
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	se thre	٠۱: دد	:	ـ ـ ـا ١			حالا حاجا:				المسلم						

Preorder trav	ersal							
First visit the p	parent and the	n visit its su	ıbtrees ir	order.				
1 DDEVED AX	TEDCAL (1)							
	CERSAL(node) ode is some	rooted tre	e node					
3 if node	$\neq NIL$ then							
	form some o			de				
	TRAVERSAL(c)							
7 end								
8 end								
Assume: ope	ration in line	4 is outputtii	ng node l	key to user				
Example								
Compute					(8)			
	241 (6)							
PRETRAVER	SAL((8))		\rightarrow	9) (3	7) (2)		
			(4)	(5)		(e)	1	
							U	
						0		

recur	sion	line		compi	utation							
level		iiie		СОПР	atation							
1		4		out	Put	8						
1		6		PRETR	AVER	SAL(9)				
a		4		outp	ut	9						
a		6	Р	RETRA	VERS	SAL((4)					
3		4		outp	sut	4						
a		6	Pi	RETRA	VERS	AL(5)					
e	ł c								6			
								9)	(3)	7)		
order c	f all out	put: 8, 9, 4	, 5, 3, 7, 2	, 6, 0, 1			(4)	(5)		(5) (1)	
							0			0		E
Posto	rder tra	versal										
First vi	sit the s	subtrees in o	order and t	hen vis	sit the	parer	t.					
2 i	nput no f node ≠	VERSAL(node) de is some NIL then th child in	rooted tr		le							
5 6 7	POST end	TRAVERSAL(child)		node							
	nd peri	verm NAMA	- Paraminan	******		=						

Exam	ple										(8)					
Comp		VERS	AL((8)				4	9)	(3 5)	3	7		3))		
recu	ırsion I		line		C	ompu	ıtatio	n									
1			5		F	POST	TRA	VERS	SAL(9)				(8		
2			5			POST	TRA	VERS	SAL(4)			9 (3	7	(a)
3			7			ou	pui	-	4				4	5			(S)
a			5			POS	TTRA	VER	SAL((5))					(0	
3			7			Оч	tpu	ł	5								
2			7		C	out t	sut	9									
و.	+ c																
order	of al	l outp	ut: 4,	5, 9,	3, 7,	0, 6,	1, 2,	8				4	9	3	8		3) (1
																⊚□	

Inorder t	ravers	al														
This orde	er is on	ly val	id for	binar	y tre	es.										
First the	left sul	otree	is vis	ited, tl	hen t	he pa	arent is	s visit	ed, a	and fin	ally t	he rig	ıht su	btree	is vis	ite
	RAVERSA				,,,,	1						8				
3 if 1	$\begin{array}{l} \mathbf{it} & node \\ node \neq N \end{array}$	IIL tl	hen		tree	node	е				9			2)		
5 Þ	TRAVEI perfor	m soi	ne op	eratio	n us	sing r	node			(4)	5		6	1		
6 IN 7 end	TRAVEI	RSAL (1	node.ri	ght)						(3	5	7	0			
								<u> </u>								
Example												8)				
Compute												٥٨				
										9				Y		
INTRAVI	EKSAL	(8)))						4				G	(1)	
									(,	3)	\ 1	5) (<u></u>			
recursior level	1	line			con	nputa	tion									
		4			NITO	A) (E.E.	2041 (
1				<u> </u>	NIK	AVER	RSAL((9))							
2		4			INTR	AVEI	RSAL	(4)))							
					. 1	1	1.									
3		5			out	put	4									
a		5			out	put	9									
		_														
2		G			INTR	AVE	RSAL	(5))							

3	4	INTF	RAVERSA	L(3))					
Ч	5	Οv	rtput	3				(8)		
1							(9)		(3)	
etc						4)	(5)	_	(6) (1) ()
							3)	7 0)	
der of all ou	tput: 4, 9, 3, 5	, 7, 8, 0, 6,	2, 1							
	, , ,	, , , ,	,							
NOTES										
	FRSAL POST	TRAVERSA	√I and IN	TRAVE	RSAL are	e all rec	cursive			
- PRETRAV	ERSAL, POST								able or	
- PRETRAV - Preorder tr	aversal is usef								able or	
- PRETRAV - Preorder tr relayed to th - Postorder	aversal is usefune children.	ul when sor	ne informa	ation ab	out the p	parent s	should I	be avail		
- PRETRAV - Preorder tr relayed to th - Postorder	raversal is usefu ne children.	ul when sor	ne informa	ation ab	out the p	parent s	should I	be avail		
- Preorder tr relayed to th - Postorder or relayed to - Inorder tra	aversal is usefune children.	ul when sor	ne informa	ation ab	bout the p	parent s	should I	be avail	/ailable	
- PRETRAV - Preorder tr relayed to th - Postorder or relayed to	raversal is usefune children. traversal is use the parent.	ul when sor	ne informa	ation ab	bout the p	parent s	should I	be avail	/ailable	
- PRETRAV - Preorder trelayed to the Postorder or relayed to - Inorder tra	raversal is usefune children. traversal is use the parent.	ul when sor	ne informa	ation ab	bout the p	parent s	should I	be avail	/ailable	
- PRETRAV - Preorder trelayed to the Postorder or relayed to - Inorder tra	raversal is usefune children. traversal is use the parent.	ul when sor	ne informa	ation ab	bout the p	parent s	should I	be avail	/ailable	
- PRETRAV - Preorder trelayed to the Postorder or relayed to - Inorder tra	raversal is usefune children. traversal is use the parent.	ul when sor	ne informa	ation ab	bout the p	parent s	should I	be avail	/ailable	
- PRETRAV - Preorder trelayed to the Postorder or relayed to - Inorder tra	raversal is usefune children. traversal is use the parent.	ul when sor	ne informa	ation ab	bout the p	parent s	should I	be avail	/ailable	
- PRETRAV - Preorder trelayed to the relayed to relayed	raversal is usefune children. traversal is use the parent.	ul when sor	ne informa	ation ab	bout the p	parent s	should I	be avail	/ailable	
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